DISCRETE MATHEMATICS, ALGORITHMS AND APPLICATIONS

Editors-in-Chief

Ding-Zhu Du University of Texas at Dallas, USA

Jinlong Shu East China Normal University, Shanghai, China



Cookies Notification

We use cookies on this site to enhance your user experience. By continuing to browse the site, you consent

to the use of our cookies. Learn More I Agree

Editorial Board

Co-Editors-in-Chief

Ding-Zhu Du University of Texas at Dallas, USA dzdu@utdallas.edu

Jinlong Shu East China Normal University, Shanghai, China jlshu@admin.ecnu.edu.cn

Advisory Editors

Tetsuo Asano (Japan Advanced Institute of Science and Technology, Japan)
Fan Chung Graham (University of California at San Diego, USA)
R.L. Graham (University of California at San Diego, USA)
D.J. Kleitman (Massachusetts Institute of Technology, Cambridge, USA)
Zhi-Ming Ma (Chinese Academy of Sciences, Beijing, China)
F.S. Roberts (Rutgers University, Piscataway, NJ, USA)
Frances Foong Yao (City University of Hong Kong, Hong Kong)

Associate Editors

Jean-Claude Bermond *(CNRS-UNSA, France)* Annalisa De Bonis *(Università degli Studi di Salerno, Italy)* Zhipeng Cai *(Georgia State University, USA)* Chiuyuan Chen *(National Chiao Tong University, Taiwan)* C.J. Colbourn *(Arizona State University, Tempe, USA)* Bhaskar Dasgupt *(University of Illinois at Chicago, USA)* M. Deza *(Ecole Normale Supérieure, Paris, France)* Zhenhua Duan (Xidian University, China) Hung-Lin Fu (National Chiaotong University, Taiwan) Suogang Gao (Hebei Normal University, Shijiazhuang, China) Xiaofeng Gao (Shanghai Jiao Tong University, China) Xiaodong Hu (Chinese Academy of Sciences, Beijing, China) Sun-Yuan Hsieh (Cheng Kung University, Taiwan) Liying Kang (Shanghai University, China) Donghyun Kim (Kennesaw State University, USA) Evangelos Kranakis (Carleton University, Canada) Deying Li (Renmin University, Beijing, China) Quan-Lin Li (Beijing University of Technology, China) Xueliang Li (Nankai University, Tianjin, China) Xiwen Lu (East China University of Science and Technology, China) Zaixin Lu (Washington State University, USA) Panos M. Pardalos (University of Florida, USA) Joseph Tonien (University of Wollongong, Australia) Alexey A. Tuzhilin (Moscow State University, Russia) Jose C. Valverde (University of Castilla-La Mancha, Spain) Wei Wang (Xi'an Jiaotong University, China) Guanghui Wang (School of Mathematics, China) Weifan Wang (Zhejiang Normal University, China) Weili Wu (The University of Texas at Dallas, USA) Dachuan Xu (Beijing University of Technology, China) Boting Yang (University of Regina, Canada) Cunquan Zhang (West Virginia University, West Virginia, USA) Xianchao Zhang (Dalian University of Technology, China) Xiaodong Zhang (Shanghai Jiao Tong University, China) Zhao Zhang (Xinjiang University, Wulumuqi, China)

Privacy policy © 2020 World Scientific Publishing Co Pte Ltd Powered by Atypon® Literatum





Home

Subject〉 Journals Books Major Reference Works Partner With Us〉 Open Access

About Us>

DISCRETE	Discrete Mathematics, Algorithms and Applications
MATHEMATICS, ALGORITHMS AND APPLICATIONS	ISSN (print): 1793-8309 ISSN (online): 1793-8317
	🔑 Tools < Share 🗳 Recommend to Library
Editors-in-Chief Ding-Chu Du Uwwenth of Touse at Dalase, USA Jinlong Shu East China Normal University, Shanghai, China	
World Scientific	
Submit an article	Subscribe

Volume 14, Issue 01 (January 2022)

Research Papers

No Access

Mixed covering arrays on graphs of small treewidth

Soumen Maity and Charles J. Colbourn

2150085

https://doi.org/10.1142/S1793830921500853

Abstract PDF/EPUB

✓ Preview Abstract

Research Papers

No Access

✓ Preview Abstract

Research Papers

No Access

On $\delta^{(k)}$ -coloring of generalized Petersen graphs

Merlin Thomas Ellumkalayil and Sudev Naduvath

2150096

https://doi.org/10.1142/S1793830921500968

Abstract PDF/EPUB

✓ Preview Abstract

Research Papers

No Access

On *r*-dynamic vertex coloring of some flower graph families

C. S. Gomathi, N. Mohanapriya, Arika Indah Kristina and Dafik

2150097

https://doi.org/10.1142/S179383092150097X

Abstract | PDF/EPUB

✓ Preview Abstract

Research Papers

No Access

Equality co-neighborhood domination in graphs

Ahmed A. Omran, Manal N. Al-Harere and Sahib Sh. Kahat

2150098

https://doi.org/10.1142/S1793830921500981

Abstract | PDF/EPUB

✓ Preview Abstract

Research Papers



On r-dynamic vertex coloring of some flower graph families

C. S. Gomathi^{*,§}, N. Mohanapriya^{*,¶}, Arika Indah Kristina^{†,‡,||} and Dafik^{†,‡,**}

*PG and Research Department of Mathematics Kongunadu Arts and Science College Tamil Nadu, India

 $^{\dagger}CGANT$ University of Jember, Indonesia

[‡]Mathematics Education Department University of Jember, Indonesia [§]gomathi9319@gmail.com ¶n.mohanamaths@gmail.com ^{||}arika.fkip@unej.ac.id **d.dafik@unej.ac.id

> Received 1 May 2020 Revised 15 November 2020 Accepted 12 January 2021 Published 3 March 2021

Let G be a simple, connected undirected graph with m vertices and n edges. Let vertex coloring c of a graph G be a mapping $c : V(G) \to S$, where |S| = k and it is k-colorable. Vertex coloring is proper if none of the any two neighboring vertices receives the similar color. An r-dynamic coloring is a proper coloring such that $|c(Nbd(v))| \ge$ $\min\{r, \deg_G(v)\}$, for each $v \in V(G)$. The r-dynamic chromatic number of a graph G is the minutest coloring k of G which is r-dynamic k-colorable and denoted by $\chi_r(G)$. By a simple view, we exhibit that $\chi_r(G) \le \chi_{r+1}(G)$, howbeit $\chi_{r+1}(G) - \chi_r(G)$ cannot be arbitrarily small. Thus, finding the result of $\chi_r(G)$ is useful. This study gave the result of r-dynamic chromatic number for the central graph, Line graph, Subdivision graph, Line of subdivision graph, Splitting graph and Mycielski graph of the Flower graph F_n denoted by $C(F_n)$, $L(F_n)$, $S(F_n)$, $L(S(F_n))$, $S(F_n)$ and $\mu(F_n)$, respectively.

Keywords: *r*-dynamic coloring; central graph; line graph; subdivision graph; line of subdivision graph; splitting graph; Mycielski graph.

Mathematics Subject Classification 2020: 05C15

1. Introduction

In this paper, the graphs $[\square]$ are considered simplistic, finite with $\delta(\min(u)(G))$, $\Delta(\max(u)(G))$ and $\chi(\operatorname{chromatic number})(G)$. Montgomery was the first person who brought out the *r*-dynamic coloring $[\square]$. An *r*-dynamic *k*-coloring is a mapping *c* from the vertex set V(G) to the set of colors $\{1, 2, \ldots, k\}$ such that (i) if $uv \in E(G)$,

 $\P Corresponding author.$

then $c(u) \neq c(v)$ and (ii) $|c(Nbd(v))| \geq \min\{r, \deg_G(v)\}$, for each $v \in V(G)$, where Nbd(v) is the set of all vertices adjacent to v, $\deg_G(v)$ is the degree of the vertex and r is the positive integer. The premier one is the adjacency condition and the paired one is the couple-adjacency condition. The r-dynamic coloring is the minimum coloring of the graph which is denoted as $\chi_r(G)$. When r = 1, then 1-dynamic chromatic number is same as the chromatic number of the graph G and if r is equal to two, it is a dynamic chromatic number. The following authors have also investigated on r-dynamic coloring [I], [Z], [I2]. Lai *et al.* [9] explained upper bounds of $\chi_r(G)$ in the pursuing lemma.

Lemma 1. $\chi_r(G) \ge \min\{r, \Delta(G)\} + 1.$

The upper bounds and lower bounds of r-dynamic chromatic number of some graphs have been explained in many research papers. Dafik *et al.* [5] studied the lower bounds of r-dynamic coloring for some graphs and gave an open problem to find the sharp lower bound for the connected graph. Alishahi [2] established that all graph G with $\chi(G) \geq 4, \chi_2(G) \leq \chi(G) + \gamma(G)$, where $\gamma(G)$ is the domination number G and also showed, for k-regular graph, $\chi(G) \geq 4, \chi_2(G) \leq \chi(G) + \alpha(G^2)$, where α is the independence number. For d-regular graph, the bounds for dynamic chromatic number in terms of independence number $\chi_d(G) \leq \chi(G) + 2\log_2\alpha(G) + 3$ were proved in [6]. In this work, the bounds of r-dynamic chromatic number are examined which gives an easier study on r-dynamic coloring.

Flower graph [3] is obtained from Helm graph H_n by joining the pendent edge to the hub vertex and it has 2n + 1 vertices and 4n edges and denoted by F_n . Figure [1] illustrate the flower graph F_4 .

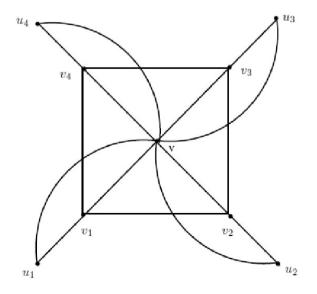


Fig. 1. Flower graph F_4 .

2150097-2

Central graph [13] is obtained from adding a new vertex to each edge at once and connects the non-adjacency vertices of the given graph.

Line graph [7] is obtained from the original graph and the vertices are adjacent if the edges of the given graph are proximity.

Subdivision graph [14] is a graph obtained by inserting a new vertex with the degree 2 for each edge of the given graph.

Line graph of Subdivision graph 14 is a graph with vertices that are acquired from the edges of the subdivision graph and the vertices are adjacent if the corresponding edges are in proximity. It is also known as **Paraline graph**.

Splitting graph S is a graph, for each vertex $V = (v_1, v_2, \ldots, v_n)$ of G, take a new vertex $V' = (v'_1, v'_2, \ldots, v'_n)$. Join each vertex of v'_i to v_j if and only if the vertex v_j is adjacent to v_i .

Mycielski graph \square is a graph which has two conditions; (i) add a vertex as given in splitting graph (ii) add another common vertex z and join the common vertex z to all points of V'.

2. Some Results on Flower Graph Families

Lemma 2. Let $C(F_n)$ be the central graph of a flower graph F_n . The lower bound for r-dynamic chromatic number of the central graph of flower graph is

$$\chi_r[C(F_n)] \ge \begin{cases} n & 1 \le r \le \delta - 1, \\ \Delta + 1 & \delta \le r \le \Delta - 1, \\ \Delta + 3 & r \ge \Delta. \end{cases}$$

Proof. $V(F_n) = \{v\} \cup \{v_i : 1 \le i \le n\} \cup \{u_i : 1 \le i \le n\}$. The order of the graph is $|V(C(F_n))| = 6n + 1$. The edge sets $\{v_i u_i, v v_i, v u_i, v_i v_{i+1} (1 \le i \le n-1), v_n v_1\}$ can be subdivided by $\{e_i, w'_i, e'_i, w_i, w_n\}$ and the size of the graph is $|E(C(F_n))| = 13n + 1$. For $1 \le r \le \delta - 1$, the vertices $V = v_i, e_i$ persuade a clique with order n in $C(F_n)$. Hence, $\chi_r[C(F_n)] \ge n$. For $\delta \le r \le \Delta - 1$ based on Lemma 1, $\chi_r(G) \ge \min\{r, \Delta(G)\} + 1$. Thus, $\chi_r[C(F_n)] \ge \min\{r, \Delta[C(F_n)]\} + 1 = \Delta + 1$. For $r \ge \Delta$ based on Lemma 1, $\chi_r[C(F_n)] \ge \min\{r, \Delta[C(F_n)]\} + 1 = \Delta + 3$. Thus, it completes the proof.

Theorem 1. For $n \ge 4$, let $C(F_n)$ be the central graph of a Flower graph F_n . Then,

$$\chi_r[C(F_n)] = \begin{cases} n & r = 1, \\ 2n+1 & \text{for } \delta = 2 \le r \le 2n-1, \\ 2n+3 & \text{for } r = \Delta = 2n, \quad n \text{ is even}, \\ 2n+4 & \text{for } r = \Delta = 2n, \quad n \text{ is odd}. \end{cases}$$

2150097-3

Proof. The *r*-dynamic coloring of $C(F_n)$ is explained in the following cases:

Case 1. When r = 1.

Based on the Lemma 2 the lower bound of $\chi_r[C(F_n)] \ge n$. The upper bound can be found out from the following coloring:

- Color the hub vertex v with color 1.
- Color the vertices v_i and u_i with *i* colors for $(1 \le i \le n)$.
- Color the vertices w_i and e_i for $(2 \le i \le n)$ with i-1 colors and w_1, e_1 with color n. Thus, 2n-2 vertices are colored.
- Color the vertices e'_i, w'_i for $(1 \le i \le n-1)$ with color n and e'_n, w'_n with color 2. Hence, couple-adjacency condition is fulfilled and therefore, $\chi_r[C(F_n)] \le n$. Hence, $\chi_r[C(F_n)] = n$.

Case 2. When $\delta = 2 \leq r \leq n - 1$.

According to Lemma 2, the lower bound of $\chi_r[C(F_n)] \ge \Delta + 1$. The *r*-dynamic (2n+1) coloring is justify to find the upper bound of $C(F_n)$:

- Assign the color *i* for the vertices v_i and e'_i for $(1 \le i \le n)$.
- Color the vertices u_i and w'_i with $n + 1, n + 2, \dots, 2n$ colors for $(1 \le i \le n)$.
- Assign the color 2n + 1 to the hub vertex v. Figure 2 illustrate the central graph of flower graph F_4 .

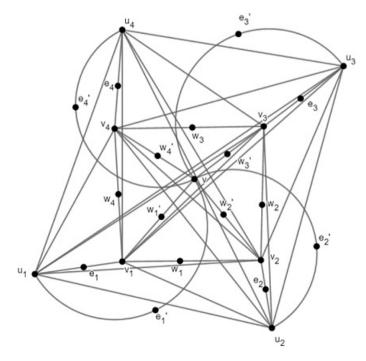


Fig. 2. Central graph of Flower graph F_4 .

2150097-4

- Color the vertex e_1 with color n and color the remaining vertices of e_i for $(2 \le i \le n)$ with $1, 2, \ldots, n-1$ colors. Thus, 2n colors are used. To satisfy the coupleadjacency coloring, we need one color. So, the coloring for the remaining n vertices is as follows:
 - (1) When n is odd, color the vertices w_i , for (i = 1, 3, ..., n-2) with 2n+1 color. To attain r-coloring, color the left over vertices w_{n-1} and w_n with color 1 and color 2.
 - (2) As n is even, color the vertices w_i for $(1 \le i \le n)$ with 2n + 1 color.

Thus, the condition of couple-adjacency coloring is satisfied and we obtain $\chi_r[C(F_n)] \leq \Delta + 1$. Hence, $\chi_r[C(F_n)] = \Delta + 1 = 2n + 1$.

Case 3. When $r = \Delta = 2n$.

Based on Lemma 2, the lower bound of $\chi_r[C(F_n)] \ge \Delta + 3$. The *r*-dynamic (2n+1) coloring is given in the following to find the upper bound of $C(F_n)$:

- Color the vertices v, u_i, w'_i, v_i and e'_i for $(1 \le i \le n)$ with the colors as given in case-2 and color the vertices e_i with 2n + 1 color for $(1 \le i \le n)$.
 - (1) As n is even, color the vertices w_i with 2n + 2 color, for (i = 1, 3, ..., n 1)and the vertices w_i with 2n + 3 color for (i = 2, 4, ..., n). Hence, $\chi_r[C(F_n)] \leq \Delta + 3$. Therefore, $\chi_r[C(F_n)] = \Delta + 3 = 2n + 3$.
 - (2) As n is odd, to satisfy the couple-adjacency condition one must add one color. So, color the vertices w_i with color 2n + 2 for (i = 1, 3, ..., n - 2) and for (i = 2, 4, 6, ..., n - 1) with color 2n + 3. Therefore, to obtain couple-adjacency condition, color the vertex w_n with 2n + 4 color. Thus, $\chi_r[C(F_n)] \leq \Delta + 4$. Hence, $\chi_r[C(F_n)] = 2n + 4$.

Lemma 3. Let $L(F_n)$ be the line graph of flower graph. The lower bound for r-dynamic chromatic number of the line graph of flower graph is

$$\chi_r[L(F_n)] \ge \begin{cases} 2n & 1 \le r \le \Delta - 3, \\ \Delta + 1 & r = \Delta - 2, \\ \Delta + 3 & \Delta - 1 \le r \le \Delta. \end{cases}$$

Proof. Let the line graph of flower graph be $L(F_n)$. The vertices of $L(F_n)$ are nothing but the edges of (F_n) ; they are $\{w_i, w'_i, e_i, e'_i\}$, for $(1 \le i \le n)$. The order of the graph $L(F_n)$ is $|V(L(F_n))| = 4n$ and the size of the graph $L(F_n)$ is $|E(L(F_n))| = 2n+3$. For $1 \le r \le \Delta - 3$, the vertices e'_i, w'_i persuade a clique of order 2n in $L(F_n)$. Thus, $\chi_r[L(F_n)] \ge 2n$. For $r = \Delta - 2$ based on Lemma 1, $\chi_r(G) \ge \min\{r, \Delta(G)\} + 1$ which implies $\chi_r[L(F_n)] \ge \min\{r, \Delta[L(F_n)]\} + 1 = \Delta + 1$. For $\Delta - 1 \le r \le \Delta$ based on Lemma 1, $\chi_r[L(F_n)] \ge \min\{r, \Delta[L(F_n)]\} + 1 = \Delta + 3$. Hence, it completes the proof.

Theorem 2. For $n \ge 5$, let $L(F_n)$ be the line graph of a Flower graph F_n . Then,

$$\chi_r[L(F_n)] = \begin{cases} 2n, & \text{for } 1 \le r \le 2n-1, \\ 2n+2, & \text{for } r = 2n \quad and \quad n \text{ is even}, \\ 2n+3, & \text{for } r = 2n \quad and \quad n \text{ is odd}, \\ 2n+5, & \text{for } 2n+1 \le r \le 2n+2 = \Delta, \\ & and \quad n \equiv 0 \pmod{6} \quad and \quad n = 5, \\ 2n+6, & \text{for } 2n+1 \le r \le 2n+2 = \Delta, \quad and \quad n \not\equiv 0 \pmod{6}. \end{cases}$$

Proof. The *r*-dynamic coloring of $L(F_n)$ is explained in the following.

Case 1. When $1 \le r \le 2n - 1$.

According to the Lemma 3, the lower bound of $\chi_r[L(F_n)] \geq 2n$. To find the upper bound color the vertices w'_i , for $(1 \leq i \leq n)$ orderly with colors i and e'_i , for $(1 \leq i \leq n)$ with the colors n + m for $(1 \leq m \leq n)$. Next, the vertex w_2 is with color n, w_1 with n - 1 color and the left over vertices of w_i for $(3 \leq i \leq n)$ are with $1, 2, \ldots, n - 2$ color. At last, the n vertices of e_i are uncolored. Hence, color the vertices e_i , for $(3 \leq i \leq n)$ with n + m colors, where $1 \leq m \leq n - 2$ and color the vertices e_1 and e_2 with the colors 2n - 1 and 2n. Thus, $\chi_r[L(F_n)] \leq 2n$. Hence, $\chi_r[L(F_n)] = 2n$. Figure \Im illustrate the line graph of flower graph F_4 .

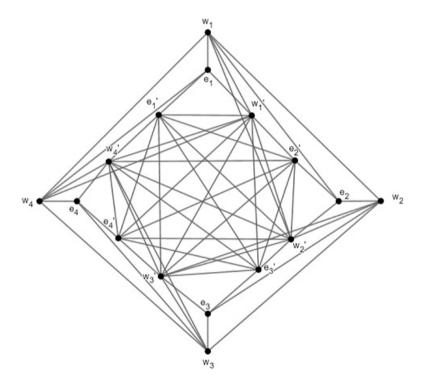


Fig. 3. Line graph of Flower graph F_4 .

Case 2. When r = 2n.

From Lemma 3, the lower bound of $\chi_r[L(F_n)] \ge \Delta + 1$. The upper bound of $L(F_n)$ can be considered from the following coloring;

- Color the vertices w'_i and e'_i as given in Case 1. Next, color the vertices w_i for $(3 \le i \le n)$ with n + m colors where, $1 \le m \le n 2$ and then color the other vertices w_2 and w_1 with 2n color and 2n 1 color.
 - (1) As n is even, the vertices e_i are with the color 2n + 1 and 2n + 2 alternatively. Hence, r-coloring has been obtained and $\chi_r[L(F_n)] \leq 2n + 2$. Therefore, $\chi_r[L(F_n)] = 2n + 2$.
 - (2) As n is odd, color the vertices e_i for $(1 \le i \le n-1)$ with 2n+1 and 2n+2 colors alternatively. In order to get the result, color the remaining vertex e_n with 2n+3 color. Hence, the couple-adjacency condition is checked out and $\chi_r[L(F_n)] \le 2n+3$. Therefore, $\chi_r[L(F_n)] = 2n+3$.

Case 3. When $2n + 1 \le r \le 2n + 2 = \triangle$.

According to Lemma 3, the lower bound of $\chi_r[L(F_n)] \ge \Delta + 3$. The upper bounds of $L(F_n)$ are demonstrated from the following way:

- The vertices w'_i and e'_i are colored in order as given in Case 1. The remaining 2n vertices of w_i and e_i remain uncolored, so color these vertices in the following way to attain the couple-adjacency condition:
 - (1) If n = 5, color the vertices w_i for $(1 \le i \le n)$ with 2n + m color where (m = 1, 2, ..., n). Next, color the vertices e_n and e_{n-1} with 2n + 1 and 2n + 2 color and color the left over vertices of e_i for $(1 \le i \le n 2)$ with 2n + m color where, (m = 3, 4, ..., n). Therefore, 2n + 5 colors are obtained and $\chi_r[L(F_n)] \le \Delta + 3$. Thus, $\chi_r[L(F_n)] = 2n + 5$.
 - (2) If $n \equiv 0 \pmod{6}$, then color the first 3t vertices of w_i with 2n + 1, 2n + 2 and 2n + 3 color in sequence, where t is the largest positive integer and $3t \leq n$. At last, remaining n vertices are uncolored. So to get (2n + 5) colors, color the vertices of e_i , with 2n + 4 and 2n + 5 color alternatively for $(1 \leq i \leq n)$. Hence, the couple-adjacency condition is satisfied and $\chi_r[L(F_n)] \leq \Delta + 3$. Thus, $\chi_r[L(F_n)] = 2n + 5$.
 - (3) In this case $n \neq 0 \pmod{6}$, assign the colors 2n + 1, 2n + 3 and 2n + 2 in sequence to the vertices w_i for $(1 \leq i \leq n-2)$. But, two vertices of w_i are uncolored. So, to obtain the result, color the vertex w_{n-1} with 2n + 4 color if n is odd else, with 2n + 5 color and color the vertex w_n with 2n + 6 color. Next, color the vertex e_1 with 2n + 2 color and the remaining vertices of e_i for $(2 \leq i \leq n-2)$ with 2n + 4 and 2n + 5 colors alternatively. Finally, two vertices of e_i are to be colored.
 - (a) the vertex e_{n-1} are colored with 2n + 2 color for $n = 7, 10, \ldots, m+1$, where m is divisible by 3 and $m \ge 6$.

- (b) Color the vertex e_{n-1} with 2n + 1 color for $n = 8, 11, \ldots, m-1$, where m is divisible by 3 and $m \ge 9$.
- (c) Color the vertex e_{n-1} with 2n+3 color for $n = 9, 15, \ldots, m+3$, where $m \equiv 0 \pmod{6}$.
- (4) If n is even, color the vertex e_n with 2n + 4 color and if n is odd then use 2n + 5 color.

Therefore, the results are obtained for the line graph of flower graph by using *r*-dynamic coloring and hence, $\chi_r[L(F_n)] \leq 2n + 6$. Therefore, $\chi_r[L(F_n)] = 2n + 6$.

Theorem 3. For $n \ge 5$, let $S(F_n)$ be the subdivision graph of a Flower graph F_n . Then,

$$\chi_r[S(F_n)] = \begin{cases} 2 & \text{for } r = 1, \\ n+1 & \text{for } \delta = 2 \le r \le n, \\ r+1 & \text{for } n+1 \le r \le \Delta = 2n. \end{cases}$$

Proof. The vertices of $S(F_n)$ are $v, v_i, u_i, w_i, w'_i, e_i$ and e'_i , for $(1 \le i \le n)$ where the vertices e_i, e'_i, w'_i are corresponding to the edge $u_i v_i, v u_i, v v_i$, the vertex w_i are corresponding to the edge $v_i v_{i+1}$ for $(1 \le i \le n-1)$ and the vertex w_n is corresponding to the edge $v_n v_1$. The order of graph $S(F_n)$ is $|V(S(F_n))| = 6n + 1$ and the size of the graph $S(F_n)$ is $|E(S(F_n))| = 7n$. The *r*-dynamic coloring of $S(F_n)$ is as follows:

Case 1. when r = 1.

Assign color 1 to the vertices v_i, u_i and v for $(1 \le i \le n)$ and the leftover vertices e_i, e'_i, w'_i and w_i for $(1 \le i \le n)$ with color 2. Hence, the couple-adjacency condition n be obtained and $\chi_r[S(F_n)] = 2$. If $\chi_r[S(F_n)] < 2$, then the couple-adjacency condition is not satisfied.

Case 2. When $\delta = 2 \leq r \leq n$.

- Assign the color *i* for the vertices v_i and e'_i where $(1 \le i \le n)$.
- Color the vertices u_i and w'_i for $(1 \le i \le n-1)$ with $2, 3, \ldots, n$ colors and the vertices u_n and w'_n with the color 1.
- The vertex e_1 is colored with color n and the vertices e_i for $(2 \le i \le n)$ with $1, 2, \ldots, n-1$ colors.
- Thus, n + 1 vertices are uncolored. So, for $(1 \le i \le n 3)$ color the vertices w_i with the color 4, 5, ..., n and the remaining vertices w_{n-2} with color 1, w_{n-1} with color 2 and w_n with color 3. To satisfy the couple-adjacency condition, the hub vertex v is colored with n+1 color. Hence, $\chi_r[S(F_n)] = n+1$. If $\chi_r[S(F_n)] < n+1$, then the couple-adjacency condition is not fulfilled.

Case 3. When $n + 1 \le r \le \Delta = 2n$.

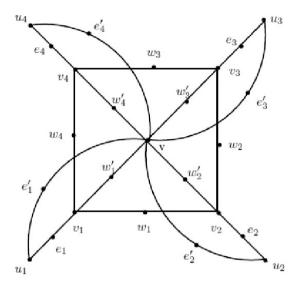


Fig. 4. Subdivision graph of Flower graph F_4 .

Figure \square illustrate the subdivision graph of flower graph F_4 . Color the vertices v, v_i, u_i, w_i, w'_i and e_i for $(1 \le i \le n)$ as given in Case 2. In order, to maintain the couple-adjacency condition, r + 1 colors are needed. When r = n + 1, assign the color n + 2 to the vertices e'_1 and e'_i for $(2 \le i \le n)$ with colors $2, 3, \ldots, n$. By continuing the coloring, for r = 2n color the vertices e'_i with $n + 2, n = 3, \ldots, 2n + 1$ colors for $(1 \le i \le n)$. Therefore, for $r \ge n$, $\chi_r[S(F_n)] = r + 1$. If $\chi_r[S(F_n)] < r + 1$, then the couple-adjacency condition is not fulfilled.

Lemma 4. Let $L[S(F_n)]$ be the paraline graph of flower graph. The lower bound for r-dynamic chromatic number of the paraline graph of flower graph is

$$\chi_r L[S(F_n)] \ge \begin{cases} 2n & 1 \le r \le \Delta - 1, \\ \Delta + 1 & r = \Delta. \end{cases}$$

Proof. The vertices of $\chi_r(L[S(F_n)])$ are a_i, b_i, c_i and d_i , which are the edges of subdivision graph $S(F_n)$. The order of the graph is $|V(L[S(F_n)])| = 8n$ and the size of the graph is $|E(L[S(F_n)])| = 2n(n+5)$. For $1 \le r \le \Delta - 1$ the vertices b_i, d_i for $i = 2, 4, \ldots, 2n$ persuade a clique of order 2n in $(L[S(F_n)])$. Thus, $\chi_r(L[S(F_n)]) \ge 2n$. For $r = \Delta$ based on Lemma 1, $\chi_r(G) \ge \min\{r, \Delta(G)\} + 1$ such that $\chi_r(L[S(F_n)]) \ge \min\{r, \Delta(L[S(F_n)])\} + 1 = \Delta + 1$. Thus, the proof is complete.

Theorem 4. For $n \ge 3$, let $L[S(F_n)]$ be the paraline graph of a Flower graph F_n . Then,

$$\chi_r(L[S(F_n)]) = \begin{cases} 2n & \text{for } 1 \le r \le 2n-1, \\ 2n+1 & \text{for } r = 2n = \Delta. \end{cases}$$

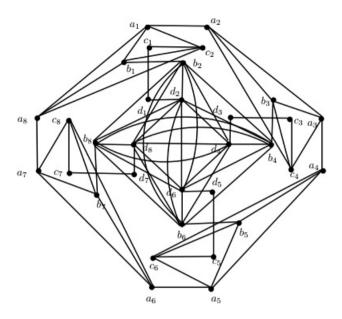


Fig. 5. Paraline graph of Flower graph F_4 .

Figure 5 illustrate the paraline graph of flower graph F_4 .

Proof. The *r*-dynamic coloring for paraline graph of flower graph is as follows:

Case 1. When $1 \le r \le 2n - 1$.

From Lemma 4, the lower bound of $\chi_r(L[S(F_n)]) \ge 2n$. However, one cannot find the sharpest lower bound. But, $\chi_r(L[S(F_n)]) = 2n$ has been considered. The upper bound can be found from coloring of $(L[S(F_n)])$ which is as follows:

- Color b_i for $i = 1, 3, \ldots, 2n-1$ with color 4. Color the vertices b_i for $i = 2, 4, \ldots, 2n$ with different colors to satisfy the couple-adjacency condition, so color the vertex b_2 with color 2, b_4 with color 1 and b_6 with color 3 and the remaining vertices of b_i for $i = 8, 10, \ldots, 2n$ with 5, $6, \ldots, n+1$ colors.
- Since, the vertices c_i are with least degree δ , color the vertices c_i for $i = 1, 3, \ldots, 2n 1$ with color 2, 1 and 3 in sequence and the other vertices c_2 , c_4 are with color 5 and c_6 are with color 6. Then, the leftover vertices c_i for $i = 8, 10, \ldots, 2n$ with color 5.
- Color the vertices d_i for i = 1, 3, ..., 2n-1 with color 3, 2 and 1 alternatively. Since the vertices d_i for i = 2, ..., 2n form a cycle (C_{n-1}) , color the vertices d_2 with color 4 and the other vertices d_i for i = 4, 6, ..., 2n with colors n+2, n+3, ..., 2n.
- At last, 2n vertices of a_i are uncolored. Hence, the coloring of the vertices a_i is as follows:
 - (1) When $n \equiv 0 \pmod{3}$, color the vertices a_i for $(1 \le i \le 2n)$ with colors 1, 2 and 3 alternatively.

- (2) When $n = 4, 7, 10, \ldots, m+1$, where $m \equiv 0 \pmod{3}$, color the vertices a_i for $(1 \le i \le 2n-2)$ with colors 1, 2 and 3 in sequence and color the vertex a_{2n-1} with color 6 and a_{2n} with color 7.
- (3) When $n = 5, 8, 11, \ldots, m-1$, where $m \equiv 0 \pmod{3}$ and $m \geq 6$, color the vertices a_i for $(1 \leq i \leq 2n-1)$ with colors 1, 2 and 3 in sequence and color the vertex a_{2n} with color 6. Hence, the couple-adjacency condition is fulfilled and $\chi_r(L[S(F_n)]) \leq 2n$. Therefore, $\chi_r(L[S(F_n)]) = 2n$.

Case 2. When r = 2n.

Based on the Lemma 4, the lower bound of $\chi_r(L[S(F_n)]) \ge \Delta + 1$. However, one cannot find the sharpest lower bound. But, $\chi_r(L[S(F_n)]) = \Delta + 1$ has been considered. To find the upper bound, color the vertices a_i, b_i and c_i for $(1 \le i \le n)$ which is given in case 1. Color d_i for $i = 1, 3, \ldots, 2n - 1$ with color 4 and the leftover vertices d_i for $i = 2, 4, \ldots, 2n$ with $n + 2, n + 3, \ldots, 2n + 1$ colors in order. Therefore, the couple-adjacency condition is satisfied and $\chi_r(L[S(F_n)]) \le \Delta + 1$ and $\chi_r(L[S(F_n)]) = 2n + 1$.

Theorem 5. For $n \ge 4$, let $S[F_n]$ be the splitting graph of Flower graph F_n . Then,

$$\chi_r(S[F_n]) = \begin{cases} 3 & \text{for } 1 \le r \le 2 = \delta, \quad n \text{ is even,} \\ 4 & \text{for } 1 \le r \le 2 = \delta, \quad n \text{ is odd,} \\ 5 & \text{for } (r \equiv 0 \mod 3) \quad and \quad r \le 5, \\ 7 & \text{for } 4 \le r \le 5, \\ r+2 & 6 \le r \le 4n = \Delta. \end{cases}$$

Proof. The vertices of $S[F_n]$ are $v, v_i, u_i, v', v'_i, u'_i$ for $(1 \le i \le n)$. The order of the graph is $|V(S[F_n])| = 4n + 2$ and the size of the graph $E(S[F_n])| = 12n$. The coloring of splitting graph of flower graph is as follows:

Case 1. When $1 \le r \le 2$.

- Color the hub vertex v and v' with color 3.
 - (1) As *n* is even, assign the color 1 and 2 to the vertices v_i and v'_i for $(1 \le i \le n)$ consecutively. Next, color the vertices u_i and u'_i with color 2 and 1 cyclically for $(1 \le i \le n)$. Hence, $\chi_r(S[F_n]) \le 3$. If $\chi_r(S[F_n]) < 3$, then the couple-adjacency condition is not fulfilled. Therefore, $\chi_r(S[F_n]) = 3$.
 - (2) When n is odd, the vertices v_i and v'_i for $(1 \le i \le n-1)$ are with color 1 and color 2 consecutively and color the vertex v_n and v'_n with color 4. Color the vertices u_i and u'_i for $(1 \le i \le n-1)$ with color 2 and color 1 consecutively and color the remaining vertex u_n and u'_n with color 1. Hence, $\chi_r(S[F_n]) \le 4$. If $\chi_r(S[F_n]) < 4$, then the couple-adjacency condition cannot be verified. Thus, $\chi_r(S[F_n]) = 4$.

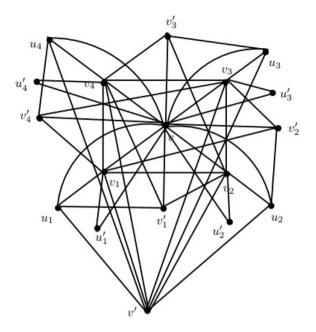


Fig. 6. Splitting graph of Flower graph F_4 .

Case 2. When $(r \equiv 0 \mod 3)$ and $r \leq 5$. Figure **6** illustrate the splitting graph of flower graph F_4 .

- Color the hub vertex v and the vertex v' with color 3 and color 4.
 - (1) When n is odd, the vertices v_i and v'_i for $(1 \le i \le n-1)$ are with color 1 and color 2 cyclically and color the vertex v_n and v'_n with color 5. Color the vertices u_i and u'_i with color 5 for $(1 \le i \le n-1)$. Next, the vertex u_n and u'_n are colored with color 1.
 - (2) When n is even, the vertices v_i and v'_i for $(1 \le i \le n)$ are with color 1 and color 2 cyclically. Color the vertices u_i and u'_i for $(1 \le i \le n)$ with color 5.

Hence, $\chi_r(S[F_n]) \leq 5$. If $\chi_r(S[F_n]) < 5$, then the couple-adjacency condition is not fulfilled. Then, $\chi_r(S[F_n]) = 5$.

Case 3. When $4 \le r \le 5$.

To satisfy the couple-adjacency condition for r = 4 one must need 7 colors. Hence, $\chi_{r=4}(S[F_n]) \leq 7$ for any *n*. If $\chi_{r=4}(S[F_n]) < 7$, the couple-adjacency condition is not satisfied. Therefore, $\chi_{r=4}(S[F_n]) = 7$. Similarly, it also satisfies the 5-dynamic coloring. Hence, $\chi_{r=4,5}(S[F_n]) = 7$ for any *n*.

Case 4. When $6 \le r \le 4n = \Delta$.

For any n, $\chi_{r=6}(S[F_n]) = 8$. If $\chi_{r=6}(S[F_n]) \leq 8$ the couple-adjacency condition is not fulfilled. Similarly, $\chi_{r=7}(S[F_n]) = 9$. Thus, one must need r+2 colors. Therefore, for any $n \chi_r(S[F_n]) = r+2$.

Theorem 6. For $n \ge 4$, let $\mu[F_n]$ be the mycielski graph of Flower graph F_n . Then,

$$\chi_r(\mu[F_n]) = \begin{cases} 4 & \text{for } 1 \le r \le 2 = \delta, \quad n \text{ is even,} \\ 5 & \text{for } 1 \le r \le 2 = \delta, \quad n \text{ is odd,} \\ r+2 & \text{for } r \equiv 0 \pmod{3} \quad and \quad n \text{ is even} \\ r+3 & \text{otherwise.} \end{cases}$$

Proof. The vertices of $\mu[F_n]$ are $z, v, v_i, u_i, v', v'_i, u'_i$ for $(1 \le i \le n)$. The order of the graph is $|V(\mu[F_n])| = 4n + 3$ and size of the graph is $|E(\mu[F_n])| = 14n + 1$. The coloring of mycielski graph of flower graph is as follows:

Case 1. When $1 \le r \le 2$.

Color the vertices u_i and u'_i for $(1 \le i \le n)$ with color 2 and color 1 orderly.

- When n is even, the vertices v_i and v'_i for $(1 \le i \le n)$ are colored with colors 1 and 2 alternatively and color the vertices u_i and u'_i for $(1 \le i \le n)$ with colors 2 and 1 in order. Color the hub vertex v and the vertex v' with color 3 and the common vertex z with color 4. Therefore, $\chi_r(\mu[F_n]) \le 4$. If $\chi_r(\mu[F_n]) < 4$, the couple-adjacency condition is not fulfilled. Hence, $\chi_r(\mu[F_n]) = 4$.
- When n is odd, the vertices v_i and v' for $(1 \le i \le n-1)$ are colored with colors 1 and 2 orderly and color the vertex v_n and v'_n with color 3. Figure 7 illustrate the Mycielski graph of flower graph F_4 . Color the hub vertex v and the vertex v' with color 4 and the common vertex z with color 5. Hence, $\chi_r(\mu[F_n]) = 5$. If $\chi_r(\mu[F_n]) < 5$, the couple-adjacency condition is not fulfilled.

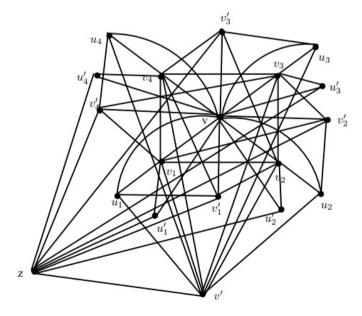


Fig. 7. Mycielski graph of Flower graph F_4 .

Case 2. When $r \equiv 0 \pmod{3}$ and n is even.

- Color the vertices v_i and v'_i with colors 1 and 2 alternatively for $(1 \le i \le n)$.
- Color the vertices u_i for $(1 \le i \le n)$ with color 5 and the vertices u'_i for $(1 \le i \le n)$ are colored with colors 2 and 1 alternatively.
- Color the hub vertex v with color 3 and the vertex v' with color 4 and the common vertex z with color 5. Thus, $\chi_r(\mu[F_n]) \leq r+2$. If $\chi_r(\mu[F_n]) < r+2$, the couple-adjacency condition is not fulfilled. Hence, $\chi_r(\mu[F_n]) = r+2$.

Case 3. otherwise.

- When $r \equiv 0 \pmod{3}$ and n is odd, color the vertices u_i and u'_i for $(1 \le i \le n)$ which is given in case 1. Color the vertices v_i for $(1 \le i \le n-1)$ with colors 1 and 2 alternatively and color the vertex v_n with color 3. Color the vertices v'_i with color 6 for $(1 \le i \le n)$. At last, color hub vertex v and assign the color 4 for the vertex v' and the common vertex z with color 5. Thus, $\chi_r(\mu[F_n]) =$ r+3. If $\chi_r(\mu[F_n]) < r+3$, the couple-adjacency condition is not verified. Hence, $\chi_r(\mu[F_n]) = r+3$.
- When $4 \le r \le 4n = \Delta$, the *r*-dynamic chromatic number for $\chi_r(\mu[F_n]) \le r+3$ for any *n*. If $\chi_r(\mu[F_n]) < r+3$, the couple-adjacency condition is not fulfilled. So, one must need r+3 colors. Hence, for any $n, \chi_r(\mu[F_n]) = r+3$.

3. Conclusion

Thus, the lower bound of the r-dynamic chromatic number of some flower graph families has been found and also some exact results have been determined. Further studies of distance graphs may give an additional insight to the r-dynamic coloring problem.

Acknowledgment

We gratefully acknowledge the support from Kongunadu Arts and Science College, Tamil Nadu, India and CGANT Research group, University of Jember, Indonesia.

References

- I. H. Agustin, D. Dafik and A. Y. Hatsya, On r-dynamic coloring of some graph operation, Int. J. Combinat. 2(1) (2016) 22–30.
- M. Alishahi, Dynamic chromatic number of regular graphs, *Discrete Appl. Math.* 160(15) (2012) 2098–2103.
- [3] R. Arundhadhi and K. Thirusangu, Star coloring of middle, total and line graph of flower graph, Int. J. Pure Appl. Math. 101(5) (2015) 691–699.
- [4] J. A. Bondy and U. S. R. Murty, Graph Theory with Applications (Macmillan Press, New York, 1976).
- [5] D. Dafik, D. E. W. Meganingtyas, K. D. Purnomo, M. D. Tarmidzi and I. H. Agustin, Several classes of graphs and their *r*-dynamic chromatic numbers, *IOP Conf. Ser. J. Phys.* 855 (2017) 012011.

- [6] A. Dehghan and A. Ahadi, Upper bounds for the 2-hued chromatic number of graphs in terms of the independence number, *Discrete Appl. Math.* 160(15) (2012) 2142– 2146.
- [7] H. Furmaczyk, J. Vernold Vivin and N. Mohanapriya, r-dynamic chromatic number of some line graphs, *Indian J. Pure Appl. Math.* 49(4) (2018) 591–600.
- [8] J. A. Gallian, A dynamic survey of graph labeling, The Electron. J. Combinat. 17 (2014) 60–62.
- [9] H. Lai, B. Montgomery and H. Poon, Upper bounds of dynamic chromatic number, Ars Combinat. 68(3) (2003) 193–201.
- [10] B. Montgomery, Dynamic coloring of graphs, ProQuest LLC, Ann Arbor, MI, Ph.D thesis, West Virginia University (2001).
- [11] N. K. Sudev, K. P. Chithra, K. A. Germina, S. Satheesh and J. Kok, On certain coloring parameters of Mycielski graphs of some graphs, *Discrete Math. Algorithms Appl.* **10**(3) (2018) 1850030.
- [12] A. Taherkhani, On r-dynamic chromatic number of graphs, Discrete Appl. Math. 201 (2016) 222–227.
- [13] J. Vernold Vivin, Harmonious coloring of total graphs, n-leaf, central graphs and circumdetic graphs, Bharathiar University, Ph.D thesis, Coimbatore, India (2007).
- [14] X. Zhang, Z. Zahid, S. Zafar, M. R. Farahani and M. F. Nadeem, Study of the para-line graphs of certain polyphenyl chains using topological indices, *Int. J. Adv. Biotechnol. Res.* 8(3) (2017) 2435–2442.