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## Research Papers

# On $r$-dynamic vertex coloring of some flower graph families 

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Let $G$ be a simple, connected undirected graph with $m$ vertices and $n$ edges. Let vertex coloring $c$ of a graph $G$ be a mapping $c: V(G) \rightarrow S$, where $|S|=k$ and it is $k$-colorable. Vertex coloring is proper if none of the any two neighboring vertices receives the similar color. An $r$-dynamic coloring is a proper coloring such that $|c(N b d(v))| \geq$ $\min \left\{r, \operatorname{deg}_{G}(v)\right\}$, for each $v \in V(G)$. The $r$-dynamic chromatic number of a graph $G$ is the minutest coloring $k$ of $G$ which is $r$-dynamic $k$-colorable and denoted by $\chi_{r}(G)$. By a simple view, we exhibit that $\chi_{r}(G) \leq \chi_{r+1}(G)$, howbeit $\chi_{r+1}(G)-\chi_{r}(G)$ cannot be arbitrarily small. Thus, finding the result of $\chi_{r}(G)$ is useful. This study gave the result of $r$-dynamic chromatic number for the central graph, Line graph, Subdivision graph, Line of subdivision graph, Splitting graph and Mycielski graph of the Flower graph $F_{n}$ denoted by $C\left(F_{n}\right), L\left(F_{n}\right), S\left(F_{n}\right), L\left(S\left(F_{n}\right)\right), S\left(F_{n}\right)$ and $\mu\left(F_{n}\right)$, respectively.

Keywords: $r$-dynamic coloring; central graph; line graph; subdivision graph; line of subdivision graph; splitting graph; Mycielski graph.

Mathematics Subject Classification 2020: 05C15

## 1. Introduction

In this paper, the graphs [4] are considered simplistic, finite with $\delta($ minimum $)(G)$, $\Delta$ (maximum) $(G)$ and $\chi$ (chromatic number) $(G)$. Montgomery was the first person who brought out the $r$-dynamic coloring [10]. An $r$-dynamic $k$-coloring is a mapping $c$ from the vertex set $V(G)$ to the set of colors $\{1,2, \ldots, k\}$ such that (i) if $u v \in E(G)$,

[^0]then $c(u) \neq c(v)$ and (ii) $|c(N b d(v))| \geq \min \left\{r, \operatorname{deg}_{G}(v)\right\}$, for each $v \in V(G)$, where $\operatorname{Nbd}(v)$ is the set of all vertices adjacent to $v, \operatorname{deg}_{G}(v)$ is the degree of the vertex and $r$ is the positive integer. The premier one is the adjacency condition and the paired one is the couple-adjacency condition. The $r$-dynamic coloring is the minimum coloring of the graph which is denoted as $\chi_{r}(G)$. When $r=1$, then 1-dynamic chromatic number is same as the chromatic number of the graph $G$ and if $r$ is equal to two, it is a dynamic chromatic number. The following authors have also investigated on $r$-dynamic coloring [1, 7, 12, Lai et al. (9] explained upper bounds of $\chi_{r}(G)$ in the pursuing lemma.

Lemma 1. $\chi_{r}(G) \geq \min \{r, \Delta(G)\}+1$.
The upper bounds and lower bounds of $r$-dynamic chromatic number of some graphs have been explained in many research papers. Dafik et al. [5] studied the lower bounds of $r$-dynamic coloring for some graphs and gave an open problem to find the sharp lower bound for the connected graph. Alishahi [2] established that all graph $G$ with $\chi(G) \geq 4, \chi_{2}(G) \leq \chi(G)+\gamma(G)$, where $\gamma(G)$ is the domination number $G$ and also showed, for $k$-regular graph, $\chi(G) \geq 4, \chi_{2}(G) \leq \chi(G)+\alpha\left(G^{2}\right)$, where $\alpha$ is the independence number. For $d$-regular graph, the bounds for dynamic chromatic number in terms of independence number $\chi_{d}(G) \leq \chi(G)+2 \log _{2} \alpha(G)+3$ were proved in [6. In this work, the bounds of $r$-dynamic chromatic number are examined which gives an easier study on $r$-dynamic coloring.

Flower graph [3] is obtained from Helm graph $H_{n}$ by joining the pendent edge to the hub vertex and it has $2 n+1$ vertices and $4 n$ edges and denoted by $F_{n}$. Figure $\square$ illustrate the flower graph $F_{4}$.


Fig. 1. Flower graph $F_{4}$.

Central graph [13] is obtained from adding a new vertex to each edge at once and connects the non-adjacency vertices of the given graph.

Line graph [7] is obtained from the original graph and the vertices are adjacent if the edges of the given graph are proximity.

Subdivision graph [14] is a graph obtained by inserting a new vertex with the degree 2 for each edge of the given graph.

Line graph of Subdivision graph [14] is a graph with vertices that are acquired from the edges of the subdivision graph and the vertices are adjacent if the corresponding edges are in proximity. It is also known as Paraline graph.

Splitting graph [8] is a graph, for each vertex $V=\left(v_{1}, v_{2}, \ldots, v_{n}\right)$ of $G$, take a new vertex $V^{\prime}=\left(v_{1}^{\prime}, v_{2}^{\prime}, \ldots, v_{n}^{\prime}\right)$. Join each vertex of $v_{i}^{\prime}$ to $v_{j}$ if and only if the vertex $v_{j}$ is adjacent to $v_{i}$.

Mycielski graph [11] is a graph which has two conditions; (i) add a vertex as given in splitting graph (ii) add another common vertex $z$ and join the common vertex $z$ to all points of $V^{\prime}$.

## 2. Some Results on Flower Graph Families

Lemma 2. Let $C\left(F_{n}\right)$ be the central graph of a flower graph $F_{n}$. The lower bound for $r$-dynamic chromatic number of the central graph of flower graph is

$$
\chi_{r}\left[C\left(F_{n}\right)\right] \geq \begin{cases}n & 1 \leq r \leq \delta-1 \\ \Delta+1 & \delta \leq r \leq \Delta-1 \\ \Delta+3 & r \geq \Delta\end{cases}
$$

Proof. $V\left(F_{n}\right)=\{v\} \cup\left\{v_{i}: 1 \leq i \leq n\right\} \cup\left\{u_{i}: 1 \leq i \leq n\right\}$. The order of the graph is $\left|V\left(C\left(F_{n}\right)\right)\right|=6 n+1$. The edge sets $\left\{v_{i} u_{i}, v v_{i}, v u_{i}, v_{i} v_{i+1}(1 \leq i \leq n-1), v_{n} v_{1}\right\}$ can be subdivided by $\left\{e_{i}, w_{i}^{\prime}, e_{i}^{\prime}, w_{i}, w_{n}\right\}$ and the size of the graph is $\left|E\left(C\left(F_{n}\right)\right)\right|=$ $13 n+1$. For $1 \leq r \leq \delta-1$, the vertices $V=v_{i}, e_{i}$ persuade a clique with order $n$ in $C\left(F_{n}\right)$. Hence, $\chi_{r}\left[C\left(F_{n}\right)\right] \geq n$. For $\delta \leq r \leq \Delta-1$ based on Lemma 1, $\chi_{r}(G) \geq$ $\min \{r, \Delta(G)\}+1$. Thus, $\chi_{r}\left[C\left(F_{n}\right)\right] \geq \min \left\{r, \Delta\left[C\left(F_{n}\right)\right]\right\}+1=\Delta+1$. For $r \geq \Delta$ based on Lemma 1, $\chi_{r}\left[C\left(F_{n}\right)\right] \geq \min \left\{r, \Delta\left[C\left(F_{n}\right)\right]\right\}+1=\Delta+3$. Thus, it completes the proof.

Theorem 1. For $n \geq 4$, let $C\left(F_{n}\right)$ be the central graph of a Flower graph $F_{n}$. Then,

$$
\chi_{r}\left[C\left(F_{n}\right)\right]= \begin{cases}n & r=1 \\ 2 n+1 & \text { for } \delta=2 \leq r \leq 2 n-1 \\ 2 n+3 & \text { for } r=\Delta=2 n, \quad n \text { is even } \\ 2 n+4 & \text { for } r=\Delta=2 n, \quad n \text { is odd }\end{cases}
$$

Proof. The $r$-dynamic coloring of $C\left(F_{n}\right)$ is explained in the following cases:
Case 1. When $r=1$.
Based on the Lemma 2 the lower bound of $\chi_{r}\left[C\left(F_{n}\right)\right] \geq n$. The upper bound can be found out from the following coloring:

- Color the hub vertex $v$ with color 1 .
- Color the vertices $v_{i}$ and $u_{i}$ with $i$ colors for $(1 \leq i \leq n)$.
- Color the vertices $w_{i}$ and $e_{i}$ for $(2 \leq i \leq n)$ with $i-1$ colors and $w_{1}, e_{1}$ with color $n$. Thus, $2 n-2$ vertices are colored.
- Color the vertices $e_{i}^{\prime}$, $w_{i}^{\prime}$ for $(1 \leq i \leq n-1)$ with color $n$ and $e_{n}^{\prime}, w_{n}^{\prime}$ with color 2. Hence, couple-adjacency condition is fulfilled and therefore, $\chi_{r}\left[C\left(F_{n}\right)\right] \leq n$. Hence, $\chi_{r}\left[C\left(F_{n}\right)\right]=n$.

Case 2. When $\delta=2 \leq r \leq n-1$.
According to Lemma 2, the lower bound of $\chi_{r}\left[C\left(F_{n}\right)\right] \geq \Delta+1$. The $r$-dynamic $(2 n+1)$ coloring is justify to find the upper bound of $C\left(F_{n}\right)$ :

- Assign the color $i$ for the vertices $v_{i}$ and $e_{i}^{\prime}$ for $(1 \leq i \leq n)$.
- Color the vertices $u_{i}$ and $w_{i}^{\prime}$ with $n+1, n+2, \ldots, 2 n$ colors for $(1 \leq i \leq n)$.
- Assign the color $2 n+1$ to the hub vertex $v$. Figure 2 illustrate the central graph of flower graph $F_{4}$.


Fig. 2. Central graph of Flower graph $F_{4}$.

- Color the vertex $e_{1}$ with color $n$ and color the remaining vertices of $e_{i}$ for $(2 \leq$ $i \leq n$ ) with $1,2, \ldots, n-1$ colors. Thus, $2 n$ colors are used. To satisfy the coupleadjacency coloring, we need one color. So, the coloring for the remaining $n$ vertices is as follows:
(1) When $n$ is odd, color the vertices $w_{i}$, for $(i=1,3, \ldots, n-2)$ with $2 n+1$ color. To attain $r$-coloring, color the left over vertices $w_{n-1}$ and $w_{n}$ with color 1 and color 2.
(2) As $n$ is even, color the vertices $w_{i}$ for $(1 \leq i \leq n)$ with $2 n+1$ color.

Thus, the condition of couple-adjacency coloring is satisfied and we obtain $\chi_{r}\left[C\left(F_{n}\right)\right] \leq \Delta+1$. Hence, $\chi_{r}\left[C\left(F_{n}\right)\right]=\Delta+1=2 n+1$.

Case 3. When $r=\Delta=2 n$.
Based on Lemma 2, the lower bound of $\chi_{r}\left[C\left(F_{n}\right)\right] \geq \Delta+3$. The $r$-dynamic $(2 n+1)$ coloring is given in the following to find the upper bound of $C\left(F_{n}\right)$ :

- Color the vertices $v, u_{i}, w_{i}^{\prime}, v_{i}$ and $e_{i}^{\prime}$ for $(1 \leq i \leq n)$ with the colors as given in case- 2 and color the vertices $e_{i}$ with $2 n+1$ color for $(1 \leq i \leq n)$.
(1) As $n$ is even, color the vertices $w_{i}$ with $2 n+2$ color, for $(i=1,3, \ldots, n-1)$ and the vertices $w_{i}$ with $2 n+3$ color for $(i=2,4, \ldots, n)$. Hence, $\chi_{r}\left[C\left(F_{n}\right)\right] \leq$ $\Delta+3$. Therefore, $\chi_{r}\left[C\left(F_{n}\right)\right]=\Delta+3=2 n+3$.
(2) As $n$ is odd, to satisfy the couple-adjacency condition one must add one color. So, color the vertices $w_{i}$ with color $2 n+2$ for $(i=1,3, \ldots, n-2)$ and for $(i=2,4,6, \ldots, n-1)$ with color $2 n+3$. Therefore, to obtain couple-adjacency condition, color the vertex $w_{n}$ with $2 n+4$ color. Thus, $\chi_{r}\left[C\left(F_{n}\right)\right] \leq \Delta+4$. Hence, $\chi_{r}\left[C\left(F_{n}\right)\right]=2 n+4$.

Lemma 3. Let $L\left(F_{n}\right)$ be the line graph of flower graph. The lower bound for $r$-dynamic chromatic number of the line graph of flower graph is

$$
\chi_{r}\left[L\left(F_{n}\right)\right] \geq \begin{cases}2 n & 1 \leq r \leq \Delta-3 \\ \Delta+1 & r=\Delta-2 \\ \Delta+3 & \Delta-1 \leq r \leq \Delta\end{cases}
$$

Proof. Let the line graph of flower graph be $L\left(F_{n}\right)$. The vertices of $L\left(F_{n}\right)$ are nothing but the edges of $\left(F_{n}\right)$; they are $\left\{w_{i}, w_{i}^{\prime}, e_{i}, e_{i}^{\prime}\right\}$, for $(1 \leq i \leq n)$. The order of the graph $L\left(F_{n}\right)$ is $\left|V\left(L\left(F_{n}\right)\right)\right|=4 n$ and the size of the graph $L\left(F_{n}\right)$ is $\left|E\left(L\left(F_{n}\right)\right)\right|=$ $2 n+3$. For $1 \leq r \leq \Delta-3$, the vertices $e_{i}^{\prime}$, $w_{i}^{\prime}$ persuade a clique of order $2 n$ in $L\left(F_{n}\right)$. Thus, $\chi_{r}\left[L\left(F_{n}\right)\right] \geq 2 n$. For $r=\Delta-2$ based on Lemma 1, $\chi_{r}(G) \geq \min \{r, \Delta(G)\}+1$ which implies $\chi_{r}\left[L\left(F_{n}\right)\right] \geq \min \left\{r, \Delta\left[L\left(F_{n}\right)\right]\right\}+1=\Delta+1$. For $\Delta-1 \leq r \leq \Delta$ based on Lemma 1, $\chi_{r}\left[L\left(F_{n}\right)\right] \geq \min \left\{r, \Delta\left[L\left(F_{n}\right)\right]\right\}+1=\Delta+3$. Hence, it completes the proof.

Theorem 2. For $n \geq 5$, let $L\left(F_{n}\right)$ be the line graph of a Flower graph $F_{n}$. Then,

$$
\chi_{r}\left[L\left(F_{n}\right)\right]= \begin{cases}2 n, & \text { for } 1 \leq r \leq 2 n-1, \\ 2 n+2, & \text { for } r=2 n \quad \text { and } n \text { is even }, \\ 2 n+3, & \text { for } r=2 n \quad \text { and } n \text { is odd, } \\ 2 n+5, & \text { for } 2 n+1 \leq r \leq 2 n+2=\triangle, \\ & \text { and } n \equiv 0 \quad(\bmod 6) \text { and } n=5, \\ 2 n+6, & \text { for } 2 n+1 \leq r \leq 2 n+2=\triangle, \quad \text { and } \quad n \not \equiv 0 \quad(\bmod 6) .\end{cases}
$$

Proof. The $r$-dynamic coloring of $L\left(F_{n}\right)$ is explained in the following.
Case 1. When $1 \leq r \leq 2 n-1$.
According to the Lemma 3, the lower bound of $\chi_{r}\left[L\left(F_{n}\right)\right] \geq 2 n$. To find the upper bound color the vertices $w_{i}^{\prime}$, for $(1 \leq i \leq n)$ orderly with colors $i$ and $e_{i}^{\prime}$, for $(1 \leq i \leq n)$ with the colors $n+m$ for $(1 \leq m \leq n)$. Next, the vertex $w_{2}$ is with color $n$, $w_{1}$ with $n-1$ color and the left over vertices of $w_{i}$ for $(3 \leq i \leq n)$ are with $1,2, \ldots, n-2$ color. At last, the $n$ vertices of $e_{i}$ are uncolored. Hence, color the vertices $e_{i}$, for $(3 \leq i \leq n)$ with $n+m$ colors, where $1 \leq m \leq n-2$ and color the vertices $e_{1}$ and $e_{2}$ with the colors $2 n-1$ and $2 n$. Thus, $\chi_{r}\left[L\left(F_{n}\right)\right] \leq 2 n$. Hence, $\chi_{r}\left[L\left(F_{n}\right)\right]=2 n$. Figure 3 illustrate the line graph of flower graph $F_{4}$.


Fig. 3. Line graph of Flower graph $F_{4}$.

Case 2. When $r=2 n$.
From Lemma 3, the lower bound of $\chi_{r}\left[L\left(F_{n}\right)\right] \geq \Delta+1$. The upper bound of $L\left(F_{n}\right)$ can be considered from the following coloring;

- Color the vertices $w_{i}^{\prime}$ and $e_{i}^{\prime}$ as given in Case 1. Next, color the vertices $w_{i}$ for $(3 \leq i \leq n)$ with $n+m$ colors where, $1 \leq m \leq n-2$ and then color the other vertices $w_{2}$ and $w_{1}$ with $2 n$ color and $2 n-1$ color.
(1) As $n$ is even, the vertices $e_{i}$ are with the color $2 n+1$ and $2 n+2$ alternatively. Hence, $r$-coloring has been obtained and $\chi_{r}\left[L\left(F_{n}\right)\right] \leq 2 n+2$. Therefore, $\chi_{r}\left[L\left(F_{n}\right)\right]=2 n+2$.
(2) As $n$ is odd, color the vertices $e_{i}$ for $(1 \leq i \leq n-1)$ with $2 n+1$ and $2 n+2$ colors alternatively. In order to get the result, color the remaining vertex $e_{n}$ with $2 n+3$ color. Hence, the couple-adjacency condition is checked out and $\chi_{r}\left[L\left(F_{n}\right)\right] \leq 2 n+3$. Therefore, $\chi_{r}\left[L\left(F_{n}\right)\right]=2 n+3$.

Case 3. When $2 n+1 \leq r \leq 2 n+2=\triangle$.
According to Lemma 3, the lower bound of $\chi_{r}\left[L\left(F_{n}\right)\right] \geq \Delta+3$. The upper bounds of $L\left(F_{n}\right)$ are demonstrated from the following way:

- The vertices $w_{i}^{\prime}$ and $e_{i}^{\prime}$ are colored in order as given in Case 1. The remaining $2 n$ vertices of $w_{i}$ and $e_{i}$ remain uncolored, so color these vertices in the following way to attain the couple-adjacency condition:
(1) If $n=5$, color the vertices $w_{i}$ for $(1 \leq i \leq n)$ with $2 n+m$ color where $(m=1,2, \ldots, n)$. Next, color the vertices $e_{n}$ and $e_{n-1}$ with $2 n+1$ and $2 n+2$ color and color the left over vertices of $e_{i}$ for $(1 \leq i \leq n-2)$ with $2 n+m$ color where, $(m=3,4, \ldots, n)$. Therefore, $2 n+5$ colors are obtained and $\chi_{r}\left[L\left(F_{n}\right)\right] \leq \Delta+3$. Thus, $\chi_{r}\left[L\left(F_{n}\right)\right]=2 n+5$.
(2) If $n \equiv 0(\bmod 6)$, then color the first $3 t$ vertices of $w_{i}$ with $2 n+1,2 n+2$ and $2 n+3$ color in sequence, where $t$ is the largest positive integer and $3 t \leq n$. At last, remaining $n$ vertices are uncolored. So to get $(2 n+5)$ colors, color the vertices of $e_{i}$, with $2 n+4$ and $2 n+5$ color alternatively for $(1 \leq i \leq n)$. Hence, the couple-adjacency condition is satisfied and $\chi_{r}\left[L\left(F_{n}\right)\right] \leq \Delta+3$. Thus, $\chi_{r}\left[L\left(F_{n}\right)\right]=2 n+5$.
(3) In this case $n \not \equiv 0(\bmod 6)$, assign the colors $2 n+1,2 n+3$ and $2 n+2$ in sequence to the vertices $w_{i}$ for $(1 \leq i \leq n-2)$. But, two vertices of $w_{i}$ are uncolored. So, to obtain the result, color the vertex $w_{n-1}$ with $2 n+4$ color if $n$ is odd else, with $2 n+5$ color and color the vertex $w_{n}$ with $2 n+6$ color. Next, color the vertex $e_{1}$ with $2 n+2$ color and the remaining vertices of $e_{i}$ for ( $2 \leq i \leq n-2$ ) with $2 n+4$ and $2 n+5$ colors alternatively. Finally, two vertices of $e_{i}$ are to be colored.
(a) the vertex $e_{n-1}$ are colored with $2 n+2$ color for $n=7,10, \ldots, m+1$, where $m$ is divisible by 3 and $m \geq 6$.
(b) Color the vertex $e_{n-1}$ with $2 n+1$ color for $n=8,11, \ldots, m-1$, where $m$ is divisible by 3 and $m \geq 9$.
(c) Color the vertex $e_{n-1}$ with $2 n+3$ color for $n=9,15, \ldots, m+3$, where $m \equiv 0(\bmod 6)$.
(4) If $n$ is even, color the vertex $e_{n}$ with $2 n+4$ color and if $n$ is odd then use $2 n+5$ color.

Therefore, the results are obtained for the line graph of flower graph by using $r$-dynamic coloring and hence, $\chi_{r}\left[L\left(F_{n}\right)\right] \leq 2 n+6$. Therefore, $\chi_{r}\left[L\left(F_{n}\right)\right]=2 n+6$.

Theorem 3. For $n \geq 5$, let $S\left(F_{n}\right)$ be the subdivision graph of a Flower graph $F_{n}$. Then,

$$
\chi_{r}\left[S\left(F_{n}\right)\right]= \begin{cases}2 & \text { for } r=1 \\ n+1 & \text { for } \delta=2 \leq r \leq n \\ r+1 & \text { for } n+1 \leq r \leq \Delta=2 n\end{cases}
$$

Proof. The vertices of $S\left(F_{n}\right)$ are $v, v_{i}, u_{i}, w_{i}, w_{i}^{\prime}, e_{i}$ and $e_{i}^{\prime}$, for $(1 \leq i \leq n)$ where the vertices $e_{i}, e_{i}^{\prime}, w_{i}^{\prime}$ are corresponding to the edge $u_{i} v_{i}, v u_{i}, v v_{i}$, the vertex $w_{i}$ are corresponding to the edge $v_{i} v_{i+1}$ for $(1 \leq i \leq n-1)$ and the vertex $w_{n}$ is corresponding to the edge $v_{n} v_{1}$. The order of graph $S\left(F_{n}\right)$ is $\left|V\left(S\left(F_{n}\right)\right)\right|=6 n+1$ and the size of the graph $S\left(F_{n}\right)$ is $\left|E\left(S\left(F_{n}\right)\right)\right|=7 n$. The $r$-dynamic coloring of $S\left(F_{n}\right)$ is as follows:

Case 1. when $r=1$.
Assign color 1 to the vertices $v_{i}, u_{i}$ and $v$ for $(1 \leq i \leq n)$ and the leftover vertices $e_{i}, e_{i}^{\prime}, w_{i}^{\prime}$ and $w_{i}$ for $(1 \leq i \leq n)$ with color 2 . Hence, the couple-adjacency condition $n$ be obtained and $\chi_{r}\left[S\left(F_{n}\right)\right]=2$. If $\chi_{r}\left[S\left(F_{n}\right)\right]<2$, then the couple-adjacency condition is not satisfied.

Case 2. When $\delta=2 \leq r \leq n$.

- Assign the color $i$ for the vertices $v_{i}$ and $e_{i}^{\prime}$ where $(1 \leq i \leq n)$.
- Color the vertices $u_{i}$ and $w_{i}^{\prime}$ for $(1 \leq i \leq n-1)$ with $2,3, \ldots, n$ colors and the vertices $u_{n}$ and $w_{n}^{\prime}$ with the color 1 .
- The vertex $e_{1}$ is colored with color $n$ and the vertices $e_{i}$ for $(2 \leq i \leq n)$ with $1,2, \ldots, n-1$ colors.
- Thus, $n+1$ vertices are uncolored. So, for $(1 \leq i \leq n-3)$ color the vertices $w_{i}$ with the color $4,5, \ldots, n$ and the remaining vertices $w_{n-2}$ with color $1, w_{n-1}$ with color 2 and $w_{n}$ with color 3 . To satisfy the couple-adjacency condition, the hub vertex $v$ is colored with $n+1$ color. Hence, $\chi_{r}\left[S\left(F_{n}\right)\right]=n+1$. If $\chi_{r}\left[S\left(F_{n}\right)\right]<n+1$, then the couple-adjacency condition is not fulfilled.

Case 3. When $n+1 \leq r \leq \Delta=2 n$.


Fig. 4. Subdivision graph of Flower graph $F_{4}$.

Figure $\mathbb{4}$ illustrate the subdivision graph of flower graph $F_{4}$. Color the vertices $v, v_{i}, u_{i}, w_{i}, w_{i}^{\prime}$ and $e_{i}$ for $(1 \leq i \leq n)$ as given in Case 2. In order, to maintain the couple-adjacency condition, $r+1$ colors are needed. When $r=n+1$, assign the color $n+2$ to the vertices $e_{1}^{\prime}$ and $e_{i}^{\prime}$ for $(2 \leq i \leq n)$ with colors $2,3, \ldots, n$. By continuing the coloring, for $r=2 n$ color the vertices $e_{i}^{\prime}$ with $n+2, n=3, \ldots, 2 n+1$ colors for $(1 \leq i \leq n)$. Therefore, for $r \geq n, \chi_{r}\left[S\left(F_{n}\right)\right]=r+1$. If $\chi_{r}\left[S\left(F_{n}\right)\right]<r+1$, then the couple-adjacency condition is not fulfilled.

Lemma 4. Let $L\left[S\left(F_{n}\right)\right]$ be the paraline graph of flower graph. The lower bound for r-dynamic chromatic number of the paraline graph of flower graph is

$$
\chi_{r} L\left[S\left(F_{n}\right)\right] \geq \begin{cases}2 n & 1 \leq r \leq \Delta-1 \\ \Delta+1 & r=\Delta\end{cases}
$$

Proof. The vertices of $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right)$ are $a_{i}, b_{i}, c_{i}$ and $d_{i}$, which are the edges of subdivision graph $S\left(F_{n}\right)$. The order of the graph is $\left|V\left(L\left[S\left(F_{n}\right)\right]\right)\right|=8 n$ and the size of the graph is $\left|E\left(L\left[S\left(F_{n}\right)\right]\right)\right|=2 n(n+5)$. For $1 \leq r \leq \Delta-1$ the vertices $b_{i}, d_{i}$ for $i=$ $2,4, \ldots, 2 n$ persuade a clique of order $2 n$ in $\left(L\left[S\left(F_{n}\right)\right]\right)$. Thus, $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right) \geq 2 n$. For $r=\Delta$ based on Lemma 1, $\chi_{r}(G) \geq \min \{r, \Delta(G)\}+1$ such that $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right) \geq$ $\min \left\{r, \Delta\left(L\left[S\left(F_{n}\right)\right]\right)\right\}+1=\Delta+1$. Thus, the proof is complete.

Theorem 4. For $n \geq 3$, let $L\left[S\left(F_{n}\right)\right]$ be the paraline graph of a Flower graph $F_{n}$. Then,

$$
\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right)= \begin{cases}2 n & \text { for } 1 \leq r \leq 2 n-1, \\ 2 n+1 & \text { for } r=2 n=\Delta\end{cases}
$$



Fig. 5. Paraline graph of Flower graph $F_{4}$.

Figure 5 illustrate the paraline graph of flower graph $F_{4}$.
Proof. The $r$-dynamic coloring for paraline graph of flower graph is as follows:
Case 1. When $1 \leq r \leq 2 n-1$.
From Lemma 4, the lower bound of $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right) \geq 2 n$. However, one cannot find the sharpest lower bound. But, $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right)=2 n$ has been considered. The upper bound can be found from coloring of $\left(L\left[S\left(F_{n}\right)\right]\right)$ which is as follows:

- Color $b_{i}$ for $i=1,3, \ldots, 2 n-1$ with color 4 . Color the vertices $b_{i}$ for $i=2,4, \ldots, 2 n$ with different colors to satisfy the couple-adjacency condition, so color the vertex $b_{2}$ with color $2, b_{4}$ with color 1 and $b_{6}$ with color 3 and the remaining vertices of $b_{i}$ for $i=8,10, \ldots, 2 n$ with $5,6, \ldots, n+1$ colors.
- Since, the vertices $c_{i}$ are with least degree $\delta$, color the vertices $c_{i}$ for $i=$ $1,3, \ldots, 2 n-1$ with color 2,1 and 3 in sequence and the other vertices $c_{2}, c_{4}$ are with color 5 and $c_{6}$ are with color 6 . Then, the leftover vertices $c_{i}$ for $i=8,10, \ldots, 2 n$ with color 5 .
- Color the vertices $d_{i}$ for $i=1,3, \ldots, 2 n-1$ with color 3,2 and 1 alternatively. Since the vertices $d_{i}$ for $i=2, \ldots, 2 n$ form a cycle $\left(C_{n-1}\right)$, color the vertices $d_{2}$ with color 4 and the other vertices $d_{i}$ for $i=4,6, \ldots, 2 n$ with colors $n+2, n+3, \ldots, 2 n$.
- At last, $2 n$ vertices of $a_{i}$ are uncolored. Hence, the coloring of the vertices $a_{i}$ is as follows:
(1) When $n \equiv 0(\bmod 3)$, color the vertices $a_{i}$ for $(1 \leq i \leq 2 n)$ with colors 1,2 and 3 alternatively.
(2) When $n=4,7,10, \ldots, m+1$, where $m \equiv 0(\bmod 3)$, color the vertices $a_{i}$ for $(1 \leq i \leq 2 n-2)$ with colors 1,2 and 3 in sequence and color the vertex $a_{2 n-1}$ with color 6 and $a_{2 n}$ with color 7 .
(3) When $n=5,8,11, \ldots, m-1$, where $m \equiv 0(\bmod 3)$ and $m \geq 6$, color the vertices $a_{i}$ for $(1 \leq i \leq 2 n-1)$ with colors 1,2 and 3 in sequence and color the vertex $a_{2 n}$ with color 6 . Hence, the couple-adjacency condition is fulfilled and $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right) \leq 2 n$. Therefore, $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right)=2 n$.

Case 2. When $r=2 n$.
Based on the Lemma 4, the lower bound of $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right) \geq \Delta+1$. However, one cannot find the sharpest lower bound. But, $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right)=\Delta+1$ has been considered. To find the upper bound, color the vertices $a_{i}, b_{i}$ and $c_{i}$ for $(1 \leq i \leq n)$ which is given in case 1 . Color $d_{i}$ for $i=1,3, \ldots, 2 n-1$ with color 4 and the leftover vertices $d_{i}$ for $i=2,4, \ldots, 2 n$ with $n+2, n+3, \ldots, 2 n+1$ colors in order. Therefore, the couple-adjacency condition is satisfied and $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right) \leq \Delta+1$ and $\chi_{r}\left(L\left[S\left(F_{n}\right)\right]\right)=2 n+1$.

Theorem 5. For $n \geq 4$, let $S\left[F_{n}\right]$ be the splitting graph of Flower graph $F_{n}$. Then,

$$
\chi_{r}\left(S\left[F_{n}\right]\right)= \begin{cases}3 & \text { for } 1 \leq r \leq 2=\delta, \quad n \text { is even } \\ 4 & \text { for } 1 \leq r \leq 2=\delta, \quad n \text { is odd } \\ 5 & \text { for }(r \equiv 0 \bmod 3) \quad \text { and } r \leq 5 \\ 7 & \text { for } 4 \leq r \leq 5 \\ r+2 & 6 \leq r \leq 4 n=\Delta\end{cases}
$$

Proof. The vertices of $S\left[F_{n}\right]$ are $v, v_{i}, u_{i}, v^{\prime}, v_{i}^{\prime}, u_{i}^{\prime}$ for $(1 \leq i \leq n)$. The order of the graph is $\left|V\left(S\left[F_{n}\right]\right)\right|=4 n+2$ and the size of the graph $E\left(S\left[F_{n}\right]\right) \mid=12 n$. The coloring of splitting graph of flower graph is as follows:

Case 1. When $1 \leq r \leq 2$.

- Color the hub vertex $v$ and $v^{\prime}$ with color 3 .
(1) As $n$ is even, assign the color 1 and 2 to the vertices $v_{i}$ and $v_{i}^{\prime}$ for $(1 \leq i \leq n)$ consecutively. Next, color the vertices $u_{i}$ and $u_{i}^{\prime}$ with color 2 and 1 cyclically for $(1 \leq i \leq n)$. Hence, $\chi_{r}\left(S\left[F_{n}\right]\right) \leq 3$. If $\chi_{r}\left(S\left[F_{n}\right]\right)<3$, then the coupleadjacency condition is not fulfilled. Therefore, $\chi_{r}\left(S\left[F_{n}\right]\right)=3$.
(2) When $n$ is odd, the vertices $v_{i}$ and $v_{i}^{\prime}$ for $(1 \leq i \leq n-1)$ are with color 1 and color 2 consecutively and color the vertex $v_{n}$ and $v_{n}^{\prime}$ with color 4 . Color the vertices $u_{i}$ and $u_{i}^{\prime}$ for $(1 \leq i \leq n-1)$ with color 2 and color 1 consecutively and color the remaining vertex $u_{n}$ and $u_{n}^{\prime}$ with color 1 . Hence, $\chi_{r}\left(S\left[F_{n}\right]\right) \leq 4$. If $\chi_{r}\left(S\left[F_{n}\right]\right)<4$, then the couple-adjacency condition cannot be verified. Thus, $\chi_{r}\left(S\left[F_{n}\right]\right)=4$.


Fig. 6. Splitting graph of Flower graph $F_{4}$.

Case 2. When $(r \equiv 0 \bmod 3)$ and $r \leq 5$. Figure 6 illustrate the spliting graph of flower graph $F_{4}$.

- Color the hub vertex $v$ and the vertex $v^{\prime}$ with color 3 and color 4 .
(1) When $n$ is odd, the vertices $v_{i}$ and $v_{i}^{\prime}$ for $(1 \leq i \leq n-1)$ are with color 1 and color 2 cyclically and color the vertex $v_{n}$ and $v_{n}^{\prime}$ with color 5 . Color the vertices $u_{i}$ and $u_{i}^{\prime}$ with color 5 for $(1 \leq i \leq n-1)$. Next, the vertex $u_{n}$ and $u_{n}^{\prime}$ are colored with color 1 .
(2) When $n$ is even, the vertices $v_{i}$ and $v_{i}^{\prime}$ for $(1 \leq i \leq n)$ are with color 1 and color 2 cyclically. Color the vertices $u_{i}$ and $u_{i}^{\prime}$ for $(1 \leq i \leq n)$ with color 5 .

Hence, $\chi_{r}\left(S\left[F_{n}\right]\right) \leq 5$. If $\chi_{r}\left(S\left[F_{n}\right]\right)<5$, then the couple-adjacency condition is not fulfilled. Then, $\chi_{r}\left(S\left[F_{n}\right]\right)=5$.
Case 3. When $4 \leq r \leq 5$.
To satisfy the couple-adjacency condition for $r=4$ one must need 7 colors. Hence, $\chi_{r=4}\left(S\left[F_{n}\right]\right) \leq 7$ for any $n$. If $\chi_{r=4}\left(S\left[F_{n}\right]\right)<7$, the couple-adjacency condition is not satisfied. Therefore, $\chi_{r=4}\left(S\left[F_{n}\right]\right)=7$. Similarly, it also satisfies the 5 -dynamic coloring. Hence, $\chi_{r=4,5}\left(S\left[F_{n}\right]\right)=7$ for any $n$.

Case 4. When $6 \leq r \leq 4 n=\Delta$.
For any $n, \chi_{r=6}\left(S\left[F_{n}\right]\right)=8$. If $\chi_{r=6}\left(S\left[F_{n}\right]\right) \leq 8$ the couple-adjacency condition is not fulfilled. Similarly, $\chi_{r=7}\left(S\left[F_{n}\right]\right)=9$. Thus, one must need $r+2$ colors. Therefore, for any $n \chi_{r}\left(S\left[F_{n}\right]\right)=r+2$.

Theorem 6. For $n \geq 4$, let $\mu\left[F_{n}\right]$ be the mycielski graph of Flower graph $F_{n}$. Then,

$$
\chi_{r}\left(\mu\left[F_{n}\right]\right)= \begin{cases}4 & \text { for } 1 \leq r \leq 2=\delta, \quad n \text { is even } \\ 5 & \text { for } 1 \leq r \leq 2=\delta, \quad n \text { is odd } \\ r+2 & \text { for } r \equiv 0 \quad(\bmod 3) \quad \text { and } n \text { is even } \\ r+3 & \text { otherwise }\end{cases}
$$

Proof. The vertices of $\mu\left[F_{n}\right]$ are $z, v, v_{i}, u_{i}, v^{\prime}, v_{i}^{\prime}, u_{i}^{\prime}$ for $(1 \leq i \leq n)$. The order of the graph is $\left|V\left(\mu\left[F_{n}\right]\right)\right|=4 n+3$ and size of the graph is $\left|E\left(\mu\left[F_{n}\right]\right)\right|=14 n+1$. The coloring of mycielski graph of flower graph is as follows:

Case 1. When $1 \leq r \leq 2$.
Color the vertices $u_{i}$ and $u_{i}^{\prime}$ for $(1 \leq i \leq n)$ with color 2 and color 1 orderly.

- When $n$ is even, the vertices $v_{i}$ and $v_{i}^{\prime}$ for $(1 \leq i \leq n)$ are colored with colors 1 and 2 alternatively and color the vertices $u_{i}$ and $u_{i}^{\prime}$ for $(1 \leq i \leq n)$ with colors 2 and 1 in order. Color the hub vertex $v$ and the vertex $v^{\prime}$ with color 3 and the common vertex $z$ with color 4 . Therefore, $\chi_{r}\left(\mu\left[F_{n}\right]\right) \leq 4$. If $\chi_{r}\left(\mu\left[F_{n}\right]\right)<4$, the couple-adjacency condition is not fulfilled. Hence, $\chi_{r}\left(\mu\left[F_{n}\right]\right)=4$.
- When $n$ is odd, the vertices $v_{i}$ and $v^{\prime}$ for $(1 \leq i \leq n-1)$ are colored with colors 1 and 2 orderly and color the vertex $v_{n}$ and $v_{n}^{\prime}$ with color 3 . Figure $\square$ illustrate the Mycielski graph of flower graph $F_{4}$. Color the hub vertex $v$ and the vertex $v^{\prime}$ with color 4 and the common vertex $z$ with color 5 . Hence, $\chi_{r}\left(\mu\left[F_{n}\right]\right)=5$. If $\chi_{r}\left(\mu\left[F_{n}\right]\right)<5$, the couple-adjacency condition is not fulfilled.


Fig. 7. Mycielski graph of Flower graph $F_{4}$.

Case 2. When $r \equiv 0(\bmod 3)$ and $n$ is even.

- Color the vertices $v_{i}$ and $v_{i}^{\prime}$ with colors 1 and 2 alternatively for $(1 \leq i \leq n)$.
- Color the vertices $u_{i}$ for $(1 \leq i \leq n)$ with color 5 and the vertices $u_{i}^{\prime}$ for $(1 \leq i \leq n)$ are colored with colors 2 and 1 alternatively.
- Color the hub vertex $v$ with color 3 and the vertex $v^{\prime}$ with color 4 and the common vertex $z$ with color 5 . Thus, $\chi_{r}\left(\mu\left[F_{n}\right]\right) \leq r+2$. If $\chi_{r}\left(\mu\left[F_{n}\right]\right)<r+2$, the couple-adjacency condition is not fulfilled. Hence, $\chi_{r}\left(\mu\left[F_{n}\right]\right)=r+2$.

Case 3. otherwise.

- When $r \equiv 0(\bmod 3)$ and $n$ is odd, color the vertices $u_{i}$ and $u_{i}^{\prime}$ for $(1 \leq i \leq n)$ which is given in case 1 . Color the vertices $v_{i}$ for $(1 \leq i \leq n-1)$ with colors 1 and 2 alternatively and color the vertex $v_{n}$ with color 3 . Color the vertices $v_{i}^{\prime}$ with color 6 for $(1 \leq i \leq n)$. At last, color hub vertex $v$ and assign the color 4 for the vertex $v^{\prime}$ and the common vertex $z$ with color 5 . Thus, $\chi_{r}\left(\mu\left[F_{n}\right]\right)=$ $r+3$. If $\chi_{r}\left(\mu\left[F_{n}\right]\right)<r+3$, the couple-adjacency condition is not verified. Hence, $\chi_{r}\left(\mu\left[F_{n}\right]\right)=r+3$.
- When $4 \leq r \leq 4 n=\Delta$, the $r$-dynamic chromatic number for $\chi_{r}\left(\mu\left[F_{n}\right]\right) \leq r+3$ for any $n$. If $\chi_{r}\left(\mu\left[F_{n}\right]\right)<r+3$, the couple-adjacency condition is not fulfilled. So, one must need $r+3$ colors. Hence, for any $n$, $\chi_{r}\left(\mu\left[F_{n}\right]\right)=r+3$.


## 3. Conclusion

Thus, the lower bound of the $r$-dynamic chromatic number of some flower graph families has been found and also some exact results have been determined. Further studies of distance graphs may give an additional insight to the $r$-dynamic coloring problem.

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